

The Depth Poset of a Filtered Lefschetz Complex

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1 Abstract

2 Taking a discrete approach to functions and dynamical systems, this paper integrates the combinatorial
3 gradients in Forman's discrete Morse theory with persistent homology to forge a unified approach to
4 function simplification. The two crucial ingredients in this effort are the *Lefschetz complex*, which
5 focuses on the homology of a cell complex at the expense of the geometry of the cells, and the
6 *shallow pairs*, which are birth-death pairs that can double as vectors in discrete Morse theory. The
7 main new concept is the *depth poset* on the birth-death pairs, which captures all simplifications
8 achieved through canceling shallow pairs. One of its linear extensions is the ordering by persistence.

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9 1 Introduction

10 The simplification of a smooth function on a manifold by canceling critical points in pairs is
11 a classic idea in Morse theory [18], whose computational execution is riddled with technical
12 difficulties. The 2- and 3-dimensional cases are of substantial practical importance in
13 geometric visualization [12], and already in three dimensions, the technical challenges abound;
14 see e.g. [11] or [16]. The purpose of this paper is to introduce new tools that help overcome
15 the technical difficulties and facilitate a clean implementation of these topological ideas.

16 We begin with an intuitive introduction of the topological idea, which we illustrate with
17 a real-valued function on a circle; see Figure 1, where A and I are identified. The graph of
18 this function may be interpreted as a mountain range in the winter, with skiers populating
19 its slopes. A skier who uses only the force of gravity can descend from a peak to one of the
20 two adjacent valleys, but this is where the journey ends. The situation improves if there is a
21 ski lift that leads up to a neighboring peak. Assuming the cost of constructing such a lift
22 increases with the height difference, we build only one lift for each peak and valley, and only
23 if it is the less expensive of the respective two choices for both, the peak and the valley; see
24 the arrows in the upper left panel of Figure 1. From the skier's viewpoint, the lifts change
25 the geometry of the mountain range as she can now reach further from most peaks. For
26 example, from the peak labeled DE, she can now reach all the way to the valleys labeled C
27 and G, but not yet beyond. The change in geometry can be visualized by leveling the peaks
28 with lifts; see the upper right panel in Figure 1 for the outcome of this operation. The new
29 geometry is reflected by the simplified graph, whose peaks and valleys are the ones without
30 lifts. We iterate the construction of lifts and this way further extend the reach of our skier
31 by simplifying the geometry; see the lower panels in Figure 1. The iteration ends with a
32 single valley and a single peak that requires no additional lifts. In the original mountain
33 range, every valley can now be reached from the remaining peak using the constructed lifts.



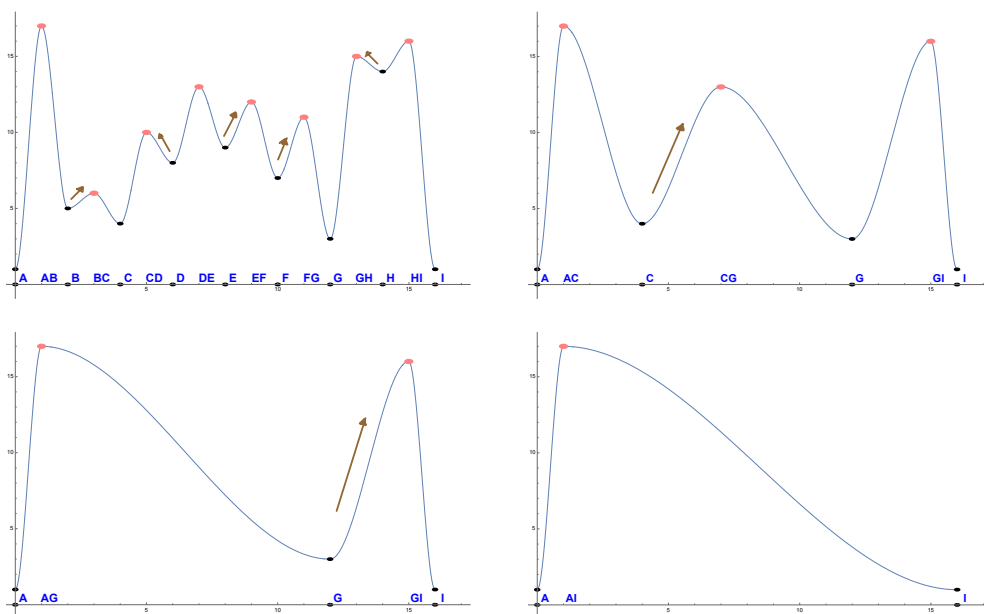
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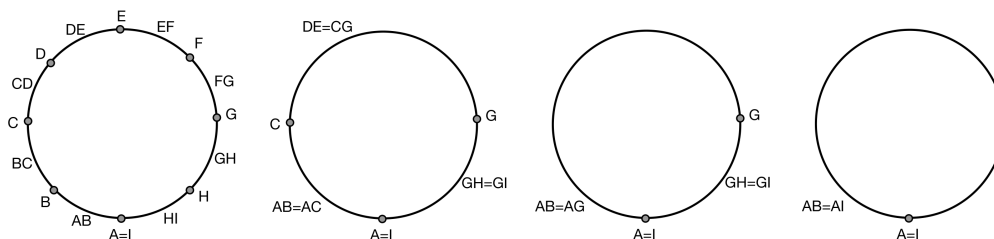
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■ Figure 1: *Upper left*: a filter over the triangulation of a circle as height function. *Upper right, lower left, and lower right*: the three derived filters obtained by canceling the shallow pairs, which are indicated by upward sloping arrows.

34 In the hope to continue the development of a combinatorial theory of dynamics started
 35 in [19, 15], our approach is fundamentally discrete and not analytic. We therefore work with
 36 a cell complex that decomposes the space, which in the above example is a decomposition of
 the circle into edges that meet in pairs at shared vertices; see Figure 2. All we need to know

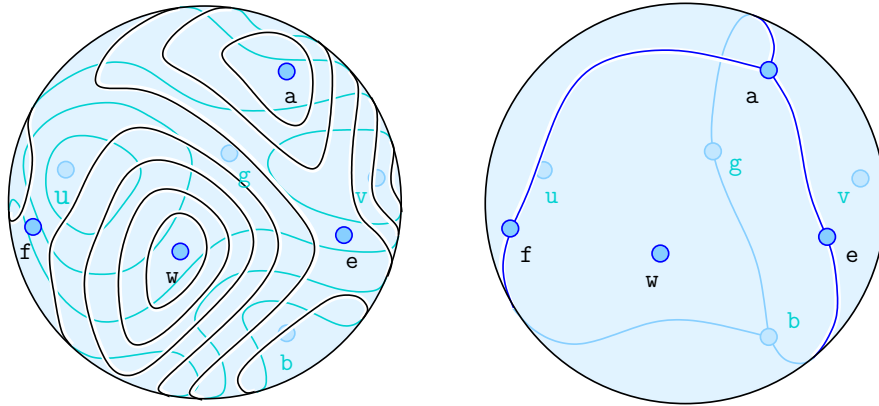


■ Figure 2: The triangulations of the circle whose vertices and edges correspond to the minima and maxima of the height functions in Figure 1. The cancellation of a peak-valley pair corresponds to removing the vertex of the valley and absorbing the edge of the peak into the neighboring edge.

37
 38 about the geometry is the height of every peak, which we store with the corresponding edge,
 39 and the height of every valley, which we store with the corresponding vertex. The leveling
 40 of a peak corresponds to absorbing the corresponding edge into the neighboring edge that
 41 shares the vertex at which the ski lift originates. As a side-effect, this vertex disappears.

42 The illustrated 1-dimensional case is of course a gross over-simplification of the general
 43 situation, so it is important to mention that the idea and our methods generalize. The
 44 2-dimensional case is illustrated in Figure 3, in which a real-valued function on the sphere
 45 on the left is represented by its critical points (minima, saddle points, maxima) and a few
 46 level lines to indicate the height differences. If we trace out the flow lines from the saddles

47 down to the minima, we decompose the sphere into the *stable manifolds* of the function,
 48 which in the generic case of a Morse–Smale function is a cell complex; see the sphere on
 the right. The geometry of this cell complex can be challenging, in particular in higher



■ Figure 3: *Left*: level sets of a height function on the sphere. There are two minima, a, b , three saddles, e, f, g , and three maxima, u, v, w . *Right*: the flow lines from the minima to the saddle points, which trace the boundaries of the maxima's influence regions.

49 dimensions. We therefore separate the geometry from the topology, which is encapsulated in
 50 the incidence relation between the cells of different dimensions. This is formalized in the
 51 notion of a *Lefschetz complex*, which we introduce in Section 2. From a purely algebraic
 52 point of view, a Lefschetz complex is merely a basis of a free chain complex. However, we
 53 feel that it is important to advocate a different point of view, with the basis elements as first
 54 class objects, and the chain complex a tool built on top of them. In our setting, a Lefschetz
 55 complex may also be viewed as a data structure for storing the homology of the complex
 56 while ignoring the geometry of its cells. Section 3 describes how this homology changes when
 57 we cancel a critical point pair and how this affects the Lefschetz complex.
 58

59 The remaining sections use these foundations to explore the simplification through
 60 successive cancellation. There is need for a global view in choosing the critical point pairs
 61 to cancel, else we may end in topologically convoluted dead-ends. We find guidance in the
 62 *vectors* of a combinatorial gradients [9], which we describe in Section 4, and in the *birth-death*
 63 *pairs* of persistent homology [7], which we explain in Section 5. The crucial concept that
 64 allows the unification of these two notions is that of a *shallow pair*, which is a sufficiently
 65 abundant type of birth-death pair that can be turned into a vector without sacrificing the
 66 acyclicity of the combinatorial vector field. By repeated cancellation of shallow pairs, we
 67 get a hierarchy of simplifications in terms of a partial order on the birth-death pairs, which
 68 we introduce in Section 6. Importantly, this poset does not depend on the order of the
 69 cancellations and represents all possible such simplifications as partitions of the poset into
 70 an upper set and a lower set.

71 We conclude this paper with a discussion of possible applications and open questions in
 72 Section 7. Among these is what originally motivated the authors of the paper to engage in
 73 this research, which is a multi-scale combinatorial theory of dynamics under local changes of
 74 the vector field. A first step would be a combinatorial Cerf theory, which restricts attention
 75 to acyclic vector fields; that is: to gradients.

76 **2 Lefschetz Complexes**

77 Introduced by the more modest name of a *complex* in the book on algebraic topology by
 78 Solomon Lefschetz [14], a Lefschetz complex may be understood as an abstraction of a cellular
 79 complex: its elements are the cells, and it stores the boundary relations between them. It is
 80 otherwise not concerned with the geometry of the cells, except that the homology of each
 81 cell is isomorphic to that of a pointed sphere of the same dimension; that is: the homology
 82 of the closed ball relative to its boundary; see the last paragraph of this section. We simplify
 83 its description by limiting ourselves to modulo-2 arithmetic.

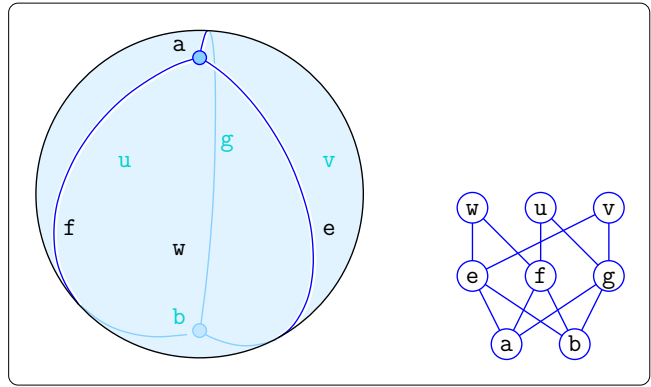
84 ► **Definition 2.1** (Lefschetz Complex). *A Lefschetz complex is a triplet (X, \dim, κ) , in which*
 85 *X is a finite set of elements called cells, $\dim: X \rightarrow \mathbb{N}_0$ maps each cell to its dimension, and*
 86 *$\kappa: X \times X \rightarrow \{0, 1\}$ is a map such that $\kappa(y, x) \neq 0$ only if $\dim y = \dim x + 1$, and*

87
$$\sum_{y \in X} \kappa(z, y) \cdot \kappa(y, x) = 0 \tag{1}$$

88 holds for all $z, x \in X$. We call x a facet of y if $\kappa(y, x) = 1$.

89 Referring to y as a p -cell if $\dim y = p$, (1) says that for every $(p + 1)$ -cell, z , and every
 90 $(p - 1)$ -cell, x , there is an even number of p -cells, y , that are facets of z and have x as a facet.
 91 We will often say that X is a Lefschetz complex, assuming \dim and κ are implicitly given.

92 As an example, consider the division of a 2-sphere into three wedges by connecting the two
 93 poles by three arcs; see Figure 4. There are two 0-cells, a, b , three 1-cells, e, f, g , and three
 94 2-cells, u, v, w , and we have $\kappa = 1$ for the pairs $(w, e), (w, f), (v, g), (v, e), (u, f), (u, g),$
 $(g, a), (g, b), (f, a), (f, b), (e, a), (e, b)$. Note that (1) is satisfied throughout.



■ Figure 4: *Left*: a division of the 2-sphere into three wedges, in which a, b are the vertices at the north-pole and south-pole, e, f, g are the arcs that connect the poles, and u, v, w are the thus created regions. *Right*: the face poset of the division.

95
 96 Since we use modulo-2 arithmetic, many of the common algebraic notions needed to
 97 define homology simplify to elementary combinatorial concepts. We write $C(X)$ for the set
 98 of subsets of X , and call $c \in C(X)$ a *chain*. It is *homogeneous* if all cells in c have the same
 99 dimension, and we call it a p -*chain* if $c \neq \emptyset$ and the common dimension of the cells in c is p .
 100 For $c, d \in C(X)$, we write $\langle c, d \rangle = \#(c \cap d)$ for the cardinality of their intersection.

101 The *boundary* of $y \in X$ is $\partial y \in C(X)$ consisting of all facets of y . Extending it linearly
 102 to chains, we get the *boundary homomorphism*, $\partial: C(X) \rightarrow C(X)$, which maps $c \subseteq X$

103 to $\partial c = \sum_{y \in c} \partial y$. Since we use modulo-2 arithmetic, the formal sum is the symmetric
 104 difference of the sets. Condition (1) guarantees $\partial\partial = 0$, so $(C(X), \partial)$ is a chain complex.
 105 $Z(X) \subseteq C(X)$ contains all *cycles*, which are the chains with zero (empty) boundary, and
 106 $B(X) \subseteq C(X)$ contains all *boundaries*, which are the chains that are the boundary of other
 107 chains. Since $\partial\partial = 0$, we have $B(X) \subseteq Z(X)$. The *homology* is the quotient of the two:
 108 $H(X) = Z(X)/B(X)$, which is the partition of $Z(X)$ into sets $B(X) + c$ with $c \in Z(X)$.
 109 This partition is well defined because $c, d \in Z(X)$ implies that $B(X) + c$ is either equal to
 110 or disjoint of $B(X) + d$. We say X is *boundaryless* if $\kappa = 0$. Then $\partial = 0$ and $H(X)$ is the
 111 partition of X into singletons.

112 The ordered pairs of cells in which the first is a facet of the second form a relation, and
 113 the transitive closure of this relation is a partial order, called the *face poset* of X . Indeed, we
 114 call x a *face* of z , and write $x \leq z$, if there is a sequence of cells $x = y_0, y_1, \dots, y_k = z$ such
 115 that y_i is a facet of y_{i+1} for $0 \leq i \leq k - 1$. The case $k = 0$ is allowed, and we call x a *proper*
 116 *face* of z if $k \geq 1$. It induces a topology on X via Alexandrov's theorem [1]: calling

- 117 ■ $U \subseteq X$ an *upper set* if $x \in U$ and $x \leq y$ implies $y \in U$;
- 118 ■ $L \subseteq X$ a *lower set* if $y \in L$ and $x \leq y$ implies $x \in L$,

119 the *open sets* in this topology are the upper sets, and the *closed sets* are the lower sets. We
 120 refer to this as the *Alexandrov topology* of X . This topology satisfies only the weakest of
 121 the separation axioms, namely that for any two distinct points, at least one has an open
 122 neighborhood that excludes the other. Indeed, if x is a proper face of y , then every open set
 123 that contains x also contains y , but not the other way round. Such a topological space is often
 124 referred to as a T_0 or Kolmogorov space. By a result of McCord [17], if the Lefschetz complex
 125 represents a regular cell complex (a cell complex with homeomorphic gluing maps), then its
 126 homology is isomorphic to the singular homology of that cell complex; see also Theorem 1.4.12
 127 on simplicial complexes and Theorem 7.1.7 on regular complexes in [3]. An example in which
 128 the Lefschetz complex does not represent a regular cell complex is illustrated in the right
 129 panel of Figure 5. Since the Lefschetz complex remembers the dimensions of the cells, we get
 130 the homology of the 2-sphere, as desired. But note that this is different from the homology
 131 of the order complex of this poset, which consists of two points.

132 Let (X, \dim, κ) be a Lefschetz complex, $Y \subseteq X$, and $\dim|_Y : Y \rightarrow \mathbb{N}_0$, $\kappa|_{Y \times Y} : Y \times Y \rightarrow$
 133 $\{0, 1\}$ the corresponding restrictions of \dim and κ . Then we call $(Y, \dim|_Y, \kappa|_{Y \times Y})$ a
 134 *Lefschetz subcomplex* of X if Y is a Lefschetz complex. It is not difficult to see that if Y is
 135 the intersection of an open set and a closed set in the Alexandrov topology of X , then Y is
 136 a Lefschetz subcomplex. A particular example is a single cell, $x \in X$, which by itself is a
 137 Lefschetz complex. Having only one cell, this Lefschetz complex is boundaryless. It follows
 138 that its homology is zero, except in dimension $p = \dim x$, in which it is $\mathbb{Z}/2\mathbb{Z}$. This is the
 139 homology of a pointed sphere of dimension p , which thus proves the earlier claim about the
 140 cells in a Lefschetz complex.

141 **3** Cancellations

142 We call an ordered pair of cells, $(s, t) \in X \times X$, a *reducible pair* in X if s is a facet of t .
 143 Given such a pair, we construct another Lefschetz complex in a process that may be viewed
 144 as a deformation retraction during which a pulls the attached cells with it to attach to the
 145 other facets of t .

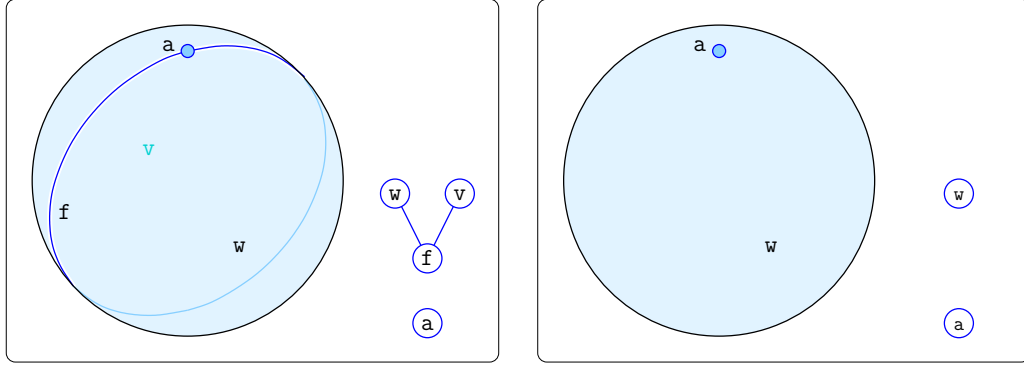
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146 ► **Definition 3.1** (Cancellation). *Given a reducible pair, (s, t) in X , set $X' = X \setminus \{s, t\}$,*
 147 *$\dim' = \dim|_{X'}$, and $\kappa': X' \times X' \rightarrow \{0, 1\}$ defined by*

$$148 \quad \kappa'(y, x) = \kappa(y, x) + \kappa(y, s) \cdot \kappa(t, x), \quad (2)$$

149 *call this operation the cancellation of (s, t) , and (X', \dim', κ') the quotient of X .*

150 In words, we increment $\kappa(y, x)$ iff s is a facet of y and x is a facet of t . See Figure 5 for
 151 examples of Lefschetz subcomplexes obtained through cancellations of reducible pairs of the Lefschetz complex in Figure 4.



152 ■ Figure 5: *Left:* the cancellations of the reducible pairs (b, e) and (g, u) simplifies the 3-division of the sphere in Figure 4 into a 2-division with a single vertex. *Right:* the further cancellation of the pair (f, v) leaves only two cells, the punctured sphere and the north-pole.

153 ► **Lemma 3.2** (Quotient). *Let (s, t) be a reducible pair in the Lefschetz complex (X, \dim, κ) .*
 154 *Then the quotient, (X', \dim', κ') , obtained by canceling (s, t) is a Lefschetz complex, and the*
 155 *corresponding boundary homomorphism is defined by mapping a cell $y \in X'$ to*

$$156 \quad \partial' y = \begin{cases} \partial y + \kappa(y, s)\partial t & \text{if } \dim' y = \dim t, \\ \partial y + \kappa(y, t)t & \text{if } \dim' y = \dim t + 1, \\ \partial y & \text{otherwise.} \end{cases} \quad (3)$$

157 The first row in (3) applies if y and t have the same dimension, and it removes the facets y
 158 shares with t and adds the other facets of t as new facets of y . The second row applies if y 's
 159 dimension exceeds that of t by one, and it removes t as a facet of y , if it was such a facet.
 160 All other cells are unaffected. We note that $\kappa(y, s) = 0$ unless $\dim' y = \dim t$, and $\kappa(y, t) = 0$
 161 unless $\dim' y = \dim t + 1$. Hence, (3) can be re-written as $\partial' y = \partial y + \kappa(y, s)\partial t + \kappa(y, t)t$.

162 **Proof.** We prove that the quotient is a Lefschetz complex by verifying that κ' satisfies (1).
 163 While κ' is formally defined only on X' , (2) makes sense also when $x = s$ and $y = t$, namely

$$164 \quad \kappa'(y, s) = \kappa(y, s) + \kappa(y, s) \cdot \kappa(t, s) = 0, \quad (4)$$

$$165 \quad \kappa'(t, x) = \kappa(t, x) + \kappa(t, s) \cdot \kappa(t, x) = 0 \quad (5)$$

166 because $\kappa(t, s) = 1$. We can therefore write the sum in (1) over all middle cells in X rather

167 than in X' :

$$168 \quad \sum_{y \in X'} \kappa'(z, y) \kappa'(y, x) = \sum_{y \in X} [\kappa(z, y) + \kappa(z, s) \kappa(t, y)] [\kappa(y, x) + \kappa(y, s) \kappa(t, x)] \quad (6)$$

$$169 \quad = \sum_{y \in X} \kappa(z, y) \kappa(y, x) + \kappa(z, s) \sum_{y \in X} \kappa(t, y) \kappa(y, x) \\ 170 \quad + \kappa(t, x) \sum_{y \in X} \kappa(z, y) \kappa(y, s) + \kappa(z, s) \kappa(t, x) \sum_{y \in X} \kappa(t, y) \kappa(y, s), \quad (7)$$

171 which vanishes because each of the four sums in (7) vanishes by assumption of X being
172 a Lefschetz complex. We omit the proof of the boundary homomorphism, which is not
173 difficult. ◀

174 We introduce three homomorphisms, which will be instrumental in proving properties of
175 the quotient. The first expresses the cancellation of (s, t) by mapping chains of X to chains
176 of X' , and the second goes the other direction, from X' to X . The homomorphisms are
177 $\pi: C(X) \rightarrow C(X')$, $\eta: C(X') \rightarrow C(X)$, and $\gamma: C(X) \rightarrow C(X)$ defined by

$$178 \quad \pi(c) = c + \langle c, s \rangle \partial t + \langle c, t \rangle t; \quad (8)$$

$$179 \quad \eta(c) = c + \langle \partial c, s \rangle t; \quad (9)$$

$$180 \quad \gamma(c) = \langle c, s \rangle t. \quad (10)$$

181 We explain (8) in words: if c contains s , then π substitutes the other facets of t for s , and if
182 c contains t , then π deletes t . To explain (9), we note that c is a chain in X' , so it is also a
183 chain in X , and ∂c denotes its boundary before the cancellation; that is: in X . If ∂c contains
184 s , then η adds t to the chain. Finally, if c contains s , then γ maps c to $\{t\}$, and else it maps
185 c to the empty chain. Inspired by work in [13], we get the following:

186 ▶ **Lemma 3.3** (Chain Homotopy). *The homomorphisms $\pi: C(X) \rightarrow C(X')$ and $\eta: C(X') \rightarrow$
187 $C(X)$ defined by a reducible pair in the Lefschetz complex, X , are chain maps, and $\gamma: C(X) \rightarrow$
188 $C(X)$ is a chain homotopy such that $\eta \circ \pi = \text{id}_{C(X)} + \partial \circ \gamma + \gamma \circ \partial$ and $\pi \circ \eta = \text{id}_{C(X')}$. In
189 particular, the chain complexes $(C(X), \partial)$ and $(C(X'), \partial')$ are chain homotopic.*

190 **Proof.** To see the first relation, we apply first π and then η to $c \in C(X)$ and rewrite the
191 terms using the two compositions of γ and ∂ :

$$192 \quad \eta \circ \pi(c) = \eta(c + \langle c, s \rangle \partial t + \langle c, t \rangle t) \quad (11)$$

$$193 \quad = c + \langle \partial c, s \rangle t + \langle c, s \rangle \partial t + \langle c, t \rangle t + \langle c, t \rangle t \quad (12)$$

$$194 \quad = \text{id}_{C(X)}(c) + \gamma \circ \partial(c) + \partial \circ \gamma(c), \quad (13)$$

195 since the last two terms in (12) cancel. To see the second relation, recall that $\pi \circ \eta$ applies
196 to chains in $C(X')$, which by construction contain neither s nor t :

$$197 \quad \pi \circ \eta(c) = \pi(c + \langle \partial c, s \rangle t) \quad (14)$$

$$198 \quad = c + \langle c, s \rangle \partial t + \langle c, t \rangle t + \langle \partial c, s \rangle t + \langle \partial c, s \rangle t \quad (15)$$

$$199 \quad = \text{id}_{C(X')}, \quad (16)$$

200 because the second and third terms in (15) vanish and the last two terms cancel. ◀

201 The existence of the chain homotopy asserted by Lemma 3.3 implies that the two Lefschetz
202 complexes, X and X' , have isomorphic homology.

203 **4 Vectors of Combinatorial Gradients**

204 Cancellations in a Lefschetz complex are not independent of each other, and one may enable
 205 or disable another. We use combinatorial gradients as introduced by Forman [9] to organize
 206 the cancellations and thus make their effect on the complex more predictable. We begin by
 207 introducing the main notions and terminology, while referring to Forman [9, 10] for a more
 208 comprehensive treatment of the background.

209 Let X be a Lefschetz complex. A *combinatorial vector field*, $V \subseteq X \times X$, is a collection
 210 of ordered pairs, called *vectors*, such that every cell belongs to at most one vector, and if
 211 $(s, t) \in V$, then s is a facet of t . Every cell that does not belong to any vector is called
 212 a *critical cell*, while the vectors are made up of *non-critical cells*. There is an associated
 213 directed graph, G_V , whose vertices are the cells in X and whose *explicit arcs* are the vectors
 214 in V . It also has *implicit arcs*, which are the pairs (y, x) such that x is a facet of y but (x, y)
 215 is not in V . A (*directed*) *path* is a sequence of vertices, x_0, x_1, \dots, x_n such that (x_i, x_{i+1})
 216 is an arc in G_V for $0 \leq i \leq n - 1$. Its *length* is n , and the path is *trivial* if $n = 0$. A path
 217 is a *cycle* if $x_0 = x_n$ and $x_i \neq x_j$ for $0 \leq i < j < n$. Since the vectors are disjoint, every
 218 explicit arc of a path is followed by an implicit arc. Moving along an explicit arc, we gain
 219 one dimension, while moving along an implicit arc, we lose a dimension. A cycle ends at the
 220 same cell it started from, which implies that it alternates between explicit and implicit arcs.

221 We call V a *combinatorial gradient* on X if G_V has only trivial cycles. A *Lyapunov*
 222 *function* for V is a map $f: X \rightarrow \mathbb{R}$, such that $f(x) = f(y)$ whenever (x, y) is an explicit arc,
 223 and $f(x) > f(y)$ whenever (x, y) is an implicit arc. It follows that $f(x_0) \geq f(x_n)$ if there is
 224 a path from x_0 to x_n .

225 **► Lemma 4.1** (Lyapunov Function). *A combinatorial vector field on a Lefschetz complex*
 226 *admits a Lyapunov function iff it is a combinatorial gradient.*

227 **Proof.** “ \implies ”: if the vector field is not a gradient, then there is at least one non-trivial cycle
 228 from a cell x_0 back to $x_n = x_0$. After every explicit arc, there is an implicit arc, so this cycle
 229 contains at least one implicit arc. But this contradicts $f(x_0) = f(x_n)$.

230 “ \impliedby ”: since G_V has no non-trivial cycle, the directed graph obtained by merging the
 231 endpoints of every explicit arc has an ordering such that x precedes y whenever (x, y) is
 232 an implicit arc. Going from left to right in this ordering, we assign a strictly decreasing
 233 sequence of function values. If a vertex corresponds to the two endpoints of an explicit arc,
 234 both endpoints get the value of the vertex. This is a Lyapunov function because $f(x) > f(y)$
 235 for every implicit arc (x, y) , and $f(x) = f(y)$ for every explicit arc (x, y) . ◀

236 An important property of a combinatorial gradient is the independence of the vectors
 237 if used in cancellations. We will see shortly, that this property crucially depends on the
 238 acyclicity of the associated digraph.

239 **► Lemma 4.2** (Independence and Acyclicity). *Let V be a combinatorial gradient on a Lefschetz*
 240 *complex, X , let (s, t) be a vector in V , and write X' for the quotient of X through canceling*
 241 *(s, t) . Then $V' = V \setminus \{s, t\}$ is a combinatorial gradient on X' .*

242 **Proof.** We first prove that V' is a combinatorial vector field on X' : if $(u, v) \neq (s, t)$ is a
 243 vector in V , then u is still a facet of v in X' . To see this, observe that at least one of
 244 $\kappa(v, s)$ and $\kappa(t, u)$ is zero, for else s, t, u, v, s would be a non-trivial cycle in G_V . Hence,
 245 $\kappa'(v, u) = \kappa(v, u)$ by (2). Since $(u, v) \in V$, we have $\kappa(v, u) = 1$, so $\kappa'(v, u) = 1$, as claimed.

246 We second show that canceling (s, t) preserves the acyclicity of the associated digraph. To
 247 derive a contradiction, assume that $G_{V'}$ has a non-trivial cycle and consider an arc (y, x) in this

248 cycle that is not arc in G_V . All explicit arcs of $G_{V'}$ are also explicit arcs of G_V , so (y, x) is an
 249 implicit arc of $G_{V'}$ and $\kappa'(y, x) = 1$ while $\kappa(y, x) = 0$. Since $\kappa'(y, x) = \kappa(y, x) + \kappa(y, s)\kappa(t, x)$
 250 by (2), this implies $\kappa(y, s) = \kappa(t, x) = 1$, so y, s, t, x is a path in G_V . By replacing all such
 251 arcs (y, x) in $G_{V'}$ by the paths y, s, t, x in G_V , we obtain a non-trivial cycle in G_V , which
 252 contradicts V being a combinatorial gradient. ◀

253 Lemma 4.2 implies that we can cancel all vectors in a combinatorial gradient, and this
 254 way obtain a Lefschetz complex in which only the critical cells remain. By Lemma 3.3, this
 255 Lefschetz complex is chain homotopic to X , and we will see shortly that it does not depend
 256 on the order in which the cancellations are applied. To this end, call a path in G_V *regular* if
 257 the vertex at which two consecutive implicit arcs meet is necessarily critical. Write $\#(y, x)$
 258 for the parity of regular paths from y to x in the associated digraph; that is: the number of
 259 such paths modulo 2.

260 ▶ **Theorem 4.3** (Morse Complex). *Let V be a combinatorial gradient on a Lefschetz complex,*
 261 *(X, \dim, κ) , and (X'', \dim'', κ'') the Lefschetz complex obtained by canceling all vectors in V .*
 262 *Then X'' is the set of critical cells of V in X , and for any two cells, $s, t \in X''$, we have*

$$263 \quad \kappa''(t, s) = \#(t, s). \quad (17)$$

264 *So (X'', \dim'', κ'') is independent of the order in which the vectors in V are cancelled.*

265 **Proof.** The only part of the theorem that still needs proof is equation (17). Let $k =$
 266 $\dim t - \dim s$, which is the surplus of implicit arcs on any path from t to s . Since t and s are
 267 critical, the first and last arcs are implicit, so the surplus is at least 1.

268 We first consider the case $k = 1$. The arcs in a path with surplus 1 alternate between
 269 implicit and explicit, which implies that every such path is regular. To prove (17), we use
 270 induction over the number of vectors in V , which we denote n . For $n = 0$, we have $X'' = X$
 271 so every arc in G_V is implicit. Hence, there is either no path from t to s or there is a path
 272 consisting of a single arc, in which case s is a facet of t . Equivalently, $\kappa''(t, s) = \#(t, s)$,
 273 which establishes the induction basis.

274 For the induction step, let V be a combinatorial gradient with $n \geq 1$ vectors and assume
 275 that (17) holds for all combinatorial gradients with $n - 1$ vectors. Letting (u, v) be a vector
 276 in V , we set $V' = V \setminus \{(u, v)\}$ and write (X', \dim', κ') for the Lefschetz complex obtained by
 277 canceling the $n - 1$ vectors in V' . By (2), we have

$$278 \quad \kappa''(t, s) = \kappa'(t, s) + \kappa'(t, u) \cdot \kappa'(v, s) \quad (18)$$

279 for all $s, t \in X''$. Writing $\#'(t, s)$ for the parity of the regular paths from t to s in $G_{V'}$, we
 280 have $\kappa'(t, s) = \#'(t, s)$, $\kappa'(t, u) = \#'(t, u)$, and $\kappa'(v, s) = \#'(v, s)$ by induction. The only
 281 difference between the associated digraphs of V and V' is the arc connecting u and v , which
 282 is explicit from u to v in G_V and implicit from v to u in $G_{V'}$. Note that a path from t to s
 283 in $G_{V'}$ necessarily avoids this implicit arc. Indeed, if it used the arc from v to u , then the
 284 preceding and succeeding arcs would also be implicit, but then the surplus of implicit arcs
 285 would be at least 2. It follows that $\#'(t, s)$ is the parity of the paths from t to s in $G_{V'}$ as
 286 well as of the paths from t to s in G_V that avoid the explicit arc from u to v . Furthermore,
 287 $\#'(t, u) \cdot \#'(v, s)$ is the parity of the paths from t to s in G_V that use this explicit arc. Hence,

$$288 \quad \#(t, s) = \#'(t, s) + \#'(t, u) \cdot \#'(v, s). \quad (19)$$

289 Comparing (18) with (19), we see that this implies $\kappa''(t, s) = \#(t, s)$, as desired. To finally
 290 prove (17) in the general case, we proceed by induction in k . The case $k = 1$ has already

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291 been established, so we fix $k \geq 2$ and assume (17) holds for all regular paths with surplus
292 less than k . We have $0 = \kappa''(t, s) = \sum_x \kappa''(t, x)\kappa''(x, s)$ by definition. We restrict ourselves
293 to vertices $x \in X''$ that satisfy $\dim x = \dim t - 1$, because otherwise $\kappa''(t, x)\kappa''(x, s) = 0$. By
294 inductive assumption, we have $\kappa''(t, x) = \#(t, x)$ and $\kappa''(x, s) = \#(x, s)$. To complete the
295 argument, we just need to ascertain that $\sum_x \#(t, x)\#(x, s) = \#(t, s)$. This is indeed the
296 case because every regular path from t to s is the concatenation of a regular path from t to
297 x and a regular path from x to s , with x being the first critical cell along the path different
298 from t . This vertex, x , is different from s because $k \geq 2$. ◀

299 Formula (17) in Theorem 4.3 shows that the quotient complex, (X'', \dim'', κ'') , is isomorphic
300 to the Morse complex as constructed in [9, Section 8]. If $\kappa''(t, s) = 1$, then this implies a
301 path from t to s in G_V such that t and s are critical and all other cells along the path are
302 non-critical. We refer to such paths as *connections*.

5 Shallow Pairs in Persistent Homology

303
304 We can simplify the Lefschetz complex beyond the Morse complex of the combinatorial
305 gradient, but for this purpose, a different algebraic structure is needed as a guide. We use
306 what we call the *depth poset* of the birth-death pairs. In a nutshell, it organizes the pairs
307 such that any linear extension of the poset gives a valid sequence of cancellations. We begin
308 with a brief introduction of persistent homology and refer to [7] for a more comprehensive
309 treatment of the background.

310 By a *filter* of a Lefschetz complex, X , we mean an injection $\phi: X \rightarrow \mathbb{R}$ such that
311 $\phi(x) < \phi(y)$ whenever x is a proper face of y . Write $X_b = \phi^{-1}(\infty, b]$ for the sublevel set
312 at $b \in \mathbb{R}$. By construction, every sublevel set of ϕ is a Lefschetz subcomplex of X . The
313 increasing sequence of distinct sublevel sets is the *filtration* induced by ϕ . To describe how
314 the homology changes as we move from one sublevel set to the next, we write $[d]_b$ for the
315 homology class of a cycle $d \in Z(X_b)$. Let $a < b$ be consecutive values of ϕ ; that is: there are
316 cells $x, y \in X$ such that $a = \phi(x)$, $b = \phi(y)$, and $X_a = X_b \setminus \{y\}$. Since ∂y is a boundary in
317 X_b , $[\partial y]_b = 0$. We say y gives death to a homology class if $[\partial y]_a \neq 0$. Otherwise, there is a
318 chain $c \in C(X_a)$ such that $\partial c = \partial y$. It follows that $c + y$ is a cycle, and $[c + y]_b \neq 0$ because
319 $c + y$ cannot be a boundary in X_b . In this case, we say y gives birth to $[c + y]_b$. We write
320 X_* and X_\times for the cells in X that give birth and death of homology classes, respectively.
321 Every cell does either, so $X_* \cap X_\times = \emptyset$ and $X_* \cup X_\times = X$. If X is a boundaryless Lefschetz
322 complex, then $X_* = X$ and $X_\times = \emptyset$.

323 Note that the homology class $[c + y]_b$ given birth to by y is not uniquely determined.
324 To fix this inconvenience, we observe that there is a unique chain $c_y \in C(X_a)$ such that
325 $\partial c_y = \partial y$ and $c_y \subseteq X_\times$. Clearly, y gives birth to the homology class of $d_y = c_y + y$. By
326 construction, $d_y \cap X_* = \{y\}$, and we call d_y the *canonical cycle* associated with y .

327 ▶ **Lemma 5.1 (Canonical Cycle Basis).** *Let $\phi: X \rightarrow \mathbb{R}$ be a filter on a Lefschetz complex, and*
328 *$t \in \mathbb{R}$ a value of ϕ . Then the canonical cycles d_y , with $y \in X_* \cap X_t$, form a basis of $Z(X_t)$.*

329 **Proof.** Let d be a cycle in X_t , let y_1, y_2, \dots, y_k be the cells in d that give birth in ϕ , denote
330 their canonical cycles by d_1, d_2, \dots, d_k , and consider $d' = d_1 + d_2 + \dots + d_k$. By construction,
331 $d' \cap X_* = d \cap X_*$. To see $d' \cap X_\times = d \cap X_\times$, note that $d + d'$ is a cycle that contains no
332 birth-giving cells. By construction, the death-giving cells do not form cycles, which leaves
333 $d + d' = 0$ as the only possibility. Hence, $d = d'$, which implies that d is a combination of the
334 canonical cycles. ◀

335 To get a basis of $H(X_t)$, we need to identify the cells $y \in X_* \cap X_t$ whose canonical
 336 cycles have not been given death by any cell in X_t yet. To do this, we construct a subset
 337 $Y_b \subseteq X_* \cap X_b$ such that the $[d_y]_b$, with $y \in Y_b$, form a basis of $H(X_b)$. The construction
 338 is inductive and paraphrases the original algorithm for computing persistent homology [8].
 339 The induction proceeds along the linear ordering of the cells induced by the filter. Letting b
 340 be the value of the first cell, y , we have $y \in X_*$ and set $Y_b = \{y\}$. For the inductive step,
 341 let $a < b$ be the values of two consecutive cells in the ordering, and let y be the second cell,
 342 so $b = \phi(y)$. If $y \in X_*$, then $Y_b = Y_a \cup \{y\}$. Otherwise, $y \in X_\times$, which implies $[\partial y]_a \neq 0$.
 343 Since Y_a defines a basis of $H(X_a)$, there is a unique subset $A \subseteq Y_a$ such that $d' = \sum_{x \in A} d_x$
 344 satisfies $[d']_a = [\partial y]_a$. We let $z \in A$ be the cell with maximum value, write $\text{birth}(y) = z$, and
 345 set $Y_b = Y_a \setminus \{\text{birth}(y)\}$. We summarize for later reference.

346 ► **Lemma 5.2 (Canonical Homology Basis).** *Let $\phi: X \rightarrow \mathbb{R}$ be a filter on a Lefschetz complex,*
 347 *and $b \in \mathbb{R}$ a value of ϕ . Then the $[d_y]_b$, with $y \in Y_b$, form a basis of $H(X_b)$.*

348 For every $y \in X_\times$, we call $(\text{birth}(y), y)$ a *birth-death pair* of ϕ , and we write $\text{BD}(\phi)$ for
 349 the collection of birth-death pairs of the filter. It is easy to see that the thus constructed
 350 map, $\text{birth}: X_\times \rightarrow X_*$, is injective. This implies that two birth-death pairs are either equal
 351 or they do not share any cell. Note however that birth is generally not bijective: cells in X_*
 352 that are not in the image represent homology classes that never die, i.e. classes in $H(X)$.

353 As an example, consider the Lefschetz complex drawn in Figure 3 on the right, with cells
 354 $X = \{\mathbf{a}, \mathbf{b}, \mathbf{e}, \mathbf{f}, \mathbf{g}, \mathbf{u}, \mathbf{v}, \mathbf{w}\}$. Assuming the filter induces the alphabetic ordering of the cells,
 355 the birth- and death-giving cells are $X_* = \{\mathbf{a}, \mathbf{b}, \mathbf{f}, \mathbf{g}, \mathbf{w}\}$ and $X_\times = \{\mathbf{e}, \mathbf{u}, \mathbf{v}\}$, respectively.
 356 Correspondingly, we have three birth-death pairs: (\mathbf{b}, \mathbf{e}) , (\mathbf{g}, \mathbf{u}) , (\mathbf{f}, \mathbf{v}) . The first two are
 357 reducible, while the third is not. This suggests we first cancel the two reducible pairs, hoping
 358 that these operations make the third pair reducible, and then cancel the third pair. This
 359 works in this particular example, but there are obstacles in the general case that require a
 360 finer distinction of the birth-death pairs, which we introduce next.

361 ► **Definition 5.3 (Shallow Pairs).** *Let $\phi: X \rightarrow \mathbb{R}$ be a filter on a Lefschetz complex. A pair*
 362 *of cells, $(s, t) \in X \times X$, is shallow if s is a facet of t , $\phi(x) \leq \phi(s)$ for all facets x of t , and*
 363 *$\phi(y) \geq \phi(t)$ for all cells y that have s as a facet.*

364 Equivalently, (s, t) is a shallow pair if s is the last facet of t in the ordering induced by the
 365 filter, and t is the first cell with facet s in this ordering. We write $\text{SH}(\phi)$ for the set of shallow
 366 pairs of the filter. Shallow pairs have been introduced in [5] under the name *apparent pairs*.
 367 They are more special than reducible birth-death pairs.

368 ► **Lemma 5.4 (Shallow Pairs are Special).** *Let $\phi: X \rightarrow \mathbb{R}$ be a filter on a Lefschetz complex.*
 369 *Every shallow pair of ϕ is a reducible birth-death pair, but not every reducible birth-death*
 370 *pair is necessarily a shallow pair of ϕ .*

371 **Proof.** We first show that there are reducible birth-death pairs that are not shallow. Consider
 372 the Lefschetz complex in Figure 3, with a filter that induces the alphabetic order except that
 373 \mathbf{v} precedes \mathbf{u} . The birth-death pairs are (\mathbf{b}, \mathbf{e}) , (\mathbf{g}, \mathbf{v}) , (\mathbf{f}, \mathbf{u}) , which are all reducible, but the
 374 third pair is not shallow because \mathbf{f} is a facet of \mathbf{u} as well as \mathbf{v} , which precedes \mathbf{u} .

375 We second prove that a shallow pair, (s, t) , is necessarily a reducible birth-death pair.
 376 Reducibility is immediate. To see that (s, t) is a birth-death pair, set $a = \phi(s)$, $b = \phi(t)$,
 377 and recall that $[\partial t]_a \neq 0$. We have $s \in \partial t$, and since it is the last cell before t in the linear
 378 ordering, $s \in Y_a$. Furthermore, s belongs to the subset $A \subseteq Y_a$ for which $d' = \sum_{x \in A} d_x$
 379 satisfies $[d']_a = [\partial t]_a$. Hence, $\text{birth}(t) = s$, as claimed. ◀

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380 A boundaryless Lefschetz complex has no shallow pair by definition. To see that every
381 other filter has at least one shallow pair, let t be the first cell in the linear ordering with
382 $\partial t \neq 0$, and let $s \in \partial t$ be the last facet in this ordering. Clearly, (s, t) is a shallow pair.

383 The remainder of this section justifies the introduction of shallow pairs by showing that
384 their cancellation does not alter the persistent homology of the filtration other than in the
385 obvious way. To state this in more technical terms, we use primes for all concepts that
386 pertain to the quotient obtained by canceling a shallow pair.

387 ► **Theorem 5.5** (Canceling a Shallow Pair). *Let $\phi: X \rightarrow \mathbb{R}$ be a filter on a Lefschetz complex,*
388 *(s, t) a shallow pair of ϕ , and $\phi': X' \rightarrow \mathbb{R}$ the filter on the quotient after canceling (s, t) .*
389 *Then $\text{SH}(\phi) \subseteq \text{SH}(\phi') \cup \{(s, t)\}$ and $\text{BD}(\phi) = \text{BD}(\phi') \cup \{(s, t)\}$.*

390 The first claim is easily established, while the second is more demanding. In the interest
391 of keeping with the flow of the current discussion, we move the proof of Theorem 5.5 to
392 Appendix A, where the theorem is restated as two claims with separate arguments.

393 Theorem 5.5 exposes a weakness of the Lefschetz complex, which may alternatively
394 be considered a strength, namely in overcoming limitations in simplifying functions on
395 non-trivial spaces reported in [2]; see also [4]. Examples are the dunce hat—which has the
396 homology of the disk but is not collapsible—and the Poincaré homology sphere—which is
397 a 3-manifold that has the homology of the ordinary 3-sphere but is not homeomorphic to
398 it. After canceling all birth-death pairs, the Lefschetz complex can no longer distinguish
399 between the disk and the dunce hat, or between the ordinary 3-sphere and the Poincaré
400 homology sphere, and this inability is crucial to cancel all birth-death pairs.

401 6 The Depth Poset

402 While Theorem 5.5 characterizes the impact of canceling a shallow pair on the persistent
403 homology, we still need to understand the impact on the Lefschetz complex. To this end,
404 we show that the collection of shallow pairs is a combinatorial gradient on the Lefschetz
405 complex, so the quotient is well defined.

406 ► **Lemma 6.1** (Shallow Pairs as Vectors). *Let $\phi: X \rightarrow \mathbb{R}$ be a filter on a Lefschetz complex.*
407 *Then the set of shallow pairs, $\text{SH}(\phi) \subseteq X \times X$, is a combinatorial gradient on X .*

408 **Proof.** We introduce $f: X \rightarrow \mathbb{R}$, which assigns the same values as the filter, except if the
409 cell is in a shallow pair, in which case it assigns the smaller of the two values:

$$410 \quad f(y) = \begin{cases} \phi(s) & \text{if } \exists (s, t) \in \text{SH}(\phi) \text{ with } y = t, \\ \phi(y) & \text{otherwise.} \end{cases} \quad (20)$$

411 According to Lemma 4.1, if f is a Lyapunov function, then $V = \text{SH}(\phi)$ is a combinatorial
412 gradient. To prove that f is indeed such a function, we consider the associated digraph, G_V .
413 Letting x be a facet of y , the two cells are either connected by an explicit arc from x to y , or
414 by an implicit arc from y to x . In the former case, (x, y) is a shallow pair of $\text{SH}(\phi)$, so we get
415 $f(x) = f(y)$ by definition of f in (20). In the latter case, we need to show that $f(x) < f(y)$.
416 Since $f(x) \leq \phi(x)$ and $\phi(x) < \phi(y)$, by definition of filter, there is something to check only
417 when $f(y) < \phi(y)$. In this case, there is a shallow pair, (s, t) , with $y = t$. Then $\phi(x) < \phi(s)$
418 because x is a facet of t but not the last one in the linear ordering induced by ϕ , which is s .
419 Since $f(x) = \phi(x)$ and $f(y) = \phi(s)$, this implies the required inequality. ◀

420 According to Lemma 6.1, we can cancel all shallow pairs, and then repeat for the quotient.
 421 By Theorem 5.5, the number of birth-death pairs decreases by the number of canceled
 422 shallow pairs. As mentioned earlier, the number of shallow pairs is strictly positive unless
 423 the Lefschetz complex is boundaryless. Hence, the number of birth-death points decreases
 424 from one iteration to the next, so the process ends after a finite number of iterations. Letting
 425 k be this number, the iteration yields a sequence of Lefschetz complexes and filters on them:

$$426 \quad \phi_j : X_j \rightarrow \mathbb{R}, \quad \text{for } 0 \leq j \leq k, \quad (21)$$

427 in which $\phi_0 : X_0 \rightarrow \mathbb{R}$ is $\phi : X \rightarrow \mathbb{R}$, and $\phi_j : X_j \rightarrow \mathbb{R}$ is the restriction of ϕ_{j-1} on the quotient
 428 of X_{j-1} obtained after canceling all shallow pairs, for $j > 0$. We call ϕ_j the j -th *derived*
 429 *filter* of ϕ . Correspondingly, we get a partition of the birth-death pairs of the initial filter
 430 into shallow pairs of the filters that arise during the iteration:

$$431 \quad \text{BD}(\phi) = \text{SH}(\phi_0) \sqcup \text{SH}(\phi_1) \sqcup \dots \sqcup \text{SH}(\phi_{k-1}). \quad (22)$$

432 This sequence is a hierarchy of simplifications of the original filter on a Lefschetz complex. A
 433 more refined hierarchy is obtained by identifying the subsets of shallow pairs that change
 434 the status of another birth-death pair from non-shallow to shallow. To define it, call a
 435 linear ordering of the birth-death pairs *cancelable* if each pair is shallow after canceling all
 436 its predecessors in the ordering. For example, every linear ordering in which all pairs in
 437 $\text{SH}(\phi_{j-1})$ precede the pairs in $\text{SH}(\phi_j)$, for $1 \leq j \leq k$, is cancelable. However, in general there
 438 are cancelable orderings that are not of this type. To cast list on them, we show that for
 439 each birth-death pair (u, v) of ϕ_j , there is a unique subset of shallow pairs of ϕ_{j-1} , such that
 440 (u, v) becomes shallow precisely at the moment all shallow pairs in this subset have been
 441 canceled. For any $S \subseteq \text{SH}(\phi_{j-1})$, we write $\phi_{j-1}^S : X_j^S \rightarrow \mathbb{R}$ for the restriction of the filter to
 442 the quotient obtained by canceling the pairs in S .

443 ► **Lemma 6.2** (Turning Shallow). *Let ϕ_{j-1} be the $(j-1)$ -st derived filter of $\phi : X \rightarrow \mathbb{R}$, and*
 444 *$(u, v) \in \text{SH}(\phi_j)$ a non-shallow birth-death pair of ϕ_{j-1} . There is a unique $S_{(u,v)} \subseteq \text{SH}(\phi_{j-1})$*
 445 *such that for every $S \subseteq \text{SH}(\phi_{j-1})$, (u, v) is a shallow pair of ϕ_{j-1}^S iff $S_{(u,v)} \subseteq S$.*

446 **Proof.** By construction of the derived filters of ϕ , (u, v) is a shallow pair of ϕ_{j-1}^S if $S =$
 447 $\text{SH}(\phi_{j-1})$; that is: when $\phi_{j-1}^S = \phi_j$. Write $V = \text{SH}(\phi_{j-1})$, and consider what this means for
 448 the paths from v to u in the associated digraph, G_V . Since $\dim v = \dim u + 1$, each such
 449 path is an alternating sequence of implicit and explicit arcs, which are shallow pairs of ϕ_{j-1} .
 450 Let $S_{(u,v)}$ be the subset of shallow pairs that belong to at least one path from v to u .

451 Let $S \subseteq V$ and recall that canceling a pair $(s, t) \in S$ corresponds to replacing the explicit
 452 arc from s to t by the implicit arc from t to s , and connecting any predecessor of s directly to
 453 any successor of t . This shortens any path that contains (s, t) by two arcs. If $S_{(u,v)} \subseteq S$, then
 454 canceling all pairs in S shortens all paths from v to u to a single implicit arc. Equivalently,
 455 (u, v) is a shallow pair of ϕ_{j-1}^S . On the other hand, if $S_{(u,v)} \not\subseteq S$, then canceling the pairs
 456 in S leaves at least one path of length at least 3 from v to u . This path together with the
 457 arc from u back to v is a non-trivial cycle, which by Lemma 6.1 implies that (u, v) is not a
 458 shallow pair of ϕ_{j-1}^S . ◀

459 We return to the collection of cancelable linear orderings of the birth-death pairs of ϕ .
 460 Each such ordering is a set of pairs, so taking the intersection is well defined.

461 ► **Definition 6.3** (Depth Poset). *The depth poset of $\phi : X \rightarrow \mathbb{R}$, denoted $\text{Depth}(\phi)$, is the*
 462 *intersection of all cancelable linear orderings of $\text{BD}(\phi)$.*

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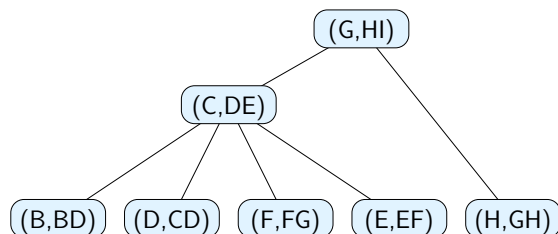
463 By definition, $\text{Depth}(\phi)$ is the largest partial order such that every cancelable linear ordering
 464 is a linear extension of this poset. We claim that it is also the smallest partial order such
 465 that every one of its linear extensions is cancelable.

466 ► **Theorem 6.4** (Cancelable Linear Orders). *Let $\phi: X \rightarrow \mathbb{R}$ be a filter on a Lefschetz complex.
 467 A linear ordering of $\text{BD}(\phi)$ is cancelable iff it is a linear extension of $\text{Depth}(\phi)$.*

468 **Proof.** By definition of the depth poset, every cancelable linear ordering of $\text{BD}(\phi)$ is a linear
 469 extension of $\text{Depth}(\phi)$. It thus suffices to prove the converse.

470 To see that every linear extension of $\text{Depth}(\phi)$ is cancelable, consider pairs $(s, t) \in$
 471 $\text{SH}(\phi_{j-1})$ and $(u, v) \in \text{SH}(\phi_j)$. By Lemma 6.2, they form a relation in the depth poset iff
 472 $(s, t) \in S_{(u,v)}$. Hence, (s, t) precedes (u, v) in every linear extension of $\text{Depth}(\phi)$. This is
 473 true for any two birth-death pairs of consecutive derived filters. It is therefore easy to show
 474 inductively that every birth-death pair is shallow after all predecessors have been canceled.
 475 Equivalently, the linear extension is cancelable. ◀

476 To give an example, we return to the graph of the 1-dimensional function in Figure 1,
 477 which we interpret as a mountain range in winter. There are eight minima and eight maxima,
 478 and since all but the global minimum and the global maximum form pairs, we have seven
 479 birth-death pairs. Figure 6 shows the depth poset of these pairs, which in this case is a tree.
 Most of the cancellation sequences of shallow pairs produce partially simplified versions of



■ Figure 6: The transitive reduction of the poset on the birth-death pairs of the function whose graph is shown in the upper left panel of Figure 1.

480 the four derived filters. The exception is when (C, DE) precedes (H, GH) . After canceling
 481 (C, DE) and before canceling (H, GH) , the graph looks like the second derived filter to the left
 482 of G and the original filter to the right of G .
 483

484 Define the *persistence* of a birth-death pair, (u, v) , as the absolute difference between
 485 their values: $\phi(v) - \phi(u)$. The persistence of a birth-death pair that is shallow during the
 486 first iteration of constructing $\text{Depth}(\phi)$ is not necessarily smaller than that of a birth-death
 487 pair that becomes shallow in later iterations. On the other hand, we will show that the pair
 488 with minimum persistence is shallow already in the first iteration. This implies that the
 489 ordering of the birth-death pairs by persistence is cancelable.

490 ► **Theorem 6.5** (Ordering by Persistence). *Let $\phi: X \rightarrow \mathbb{R}$ be a filter on a Lefschetz. Then
 491 the ordering of the birth-death pairs by persistence is a linear extension of $\text{Depth}(\phi)$.*

492 **Proof.** Assuming the minimum persistence birth-death pair is shallow, we can cancel it
 493 and iterate. This way we cancel the birth-death pairs in the order of their persistence, and
 494 since every cancelable ordering is a linear extension of the depth poset, $\text{Depth}(\phi)$, so is the
 495 ordering by persistence.

496 To see that a minimum persistence birth-death pair is shallow we prove the contraposition;
 497 that is: a birth-death pair that is not shallow does not minimize persistence. Suppose

498 $(s, t) \in \text{BD}(\phi)$ is not shallow, and fix a cancelable linear ordering in which (s, t) comes
 499 immediately after the pair $(u, v) \in \text{BD}(\phi)$, whose cancellation changes the status of (s, t)
 500 from non-shallow to shallow. Let X be the cells before canceling (u, v) , and set $V = \{(u, v)\}$,
 501 so the associated digraph, G_V has a single explicit arc. The pair (s, t) may or may not be
 502 reducible before canceling (u, v) . In the latter case, there is no arc connecting the two nodes,
 503 and because (s, t) is reducible afterwards, we know from (17) that there is a path from t to
 504 s in G_V . Since there is only one explicit arc in G_V , the path must be t, u, v, s . Since (u, v)
 505 is shallow at the time it is canceled, we have $\phi(s) < \phi(u) < \phi(v) < \phi(t)$, which shows that
 506 (u, v) is a birth-death pair with smaller persistence than (s, t) .

507 There remains the case when (s, t) is reducible before canceling (u, v) . Since (s, t) is not
 508 yet shallow, there is a cell, y , with $\phi(y) < \phi(t)$ that has s as a facet, or there is a facet, w ,
 509 of t with $\phi(s) < \phi(w)$. The two cases are symmetric. We therefore consider only the latter
 510 and assume that w is the last facet of t in the fixed linear ordering. We know that (s, t) gets
 511 shallow eventually, so w must be canceled prior to that event. Since w is the last facet of t , it
 512 gives birth, so there exists a cell z such that (w, z) is a birth-death pair that becomes shallow
 513 before (s, t) . Hence, $\phi(s) < \phi(w) < \phi(z) < \phi(t)$, which implies that (w, z) is a birth-death
 514 pair with strictly smaller persistence than (s, t) . ◀

515 7 Discussion

516 This paper introduces tools for the study of the dynamics under changing vector fields from
 517 a combinatorial viewpoint. Starting with the simpler case of a combinatorial gradient, it
 518 would be interesting to develop a combinatorial Cerf theory that classifies the non-generic
 519 critical cases, which necessarily arise when a vector field changes continuously. Perhaps the
 520 non-generic cases require multi-vectors consisting of more than two cells, which would go
 521 beyond the theory as introduced in [9]. This general topic is related to computing vineyards,
 522 as studied in [6]. At this time, the details of this relationship are unclear, primarily because
 523 of the different constraints imposed by the data structures representing the data, which are
 524 Lefschetz complexes in this paper and simplicial complexes in [6].

525 When we go to continuous vector fields more general than gradients, we observe recurring
 526 patterns that are more complicated than critical points, such as attractive or repulsive closed
 527 curves and more. The study of such phenomena may benefit from the ability to cancel a cell
 528 with one of its facets even if this creates a cell whose homology is different from that of a
 529 pointed sphere. Similarly, in the simplification of a smooth map, the restriction to cells with
 530 simple homology seems artificial and can sometimes be inconvenient.

531 We finally address the algorithmic aspects of the work described in this paper. From a
 532 computational point of view, the Lefschetz complex is an abstract data type that supports
 533 the cancellation of a reducible pair of cells. The repeated application of this operation
 534 shortens paths in the associated digraph, and the original paths between critical cells can
 535 be recovered by following the cancellations backward in time. All these operations reduce
 536 to the manipulation of lists and graphs, which are likely to have very fast implementations.
 537 It would be worthwhile to develop the details of these algorithms and to experiment with
 538 different data structures implementing them.

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A Proof of Theorem 5.5

580

581 Theorem 5.5 makes two claims, one about the shallow pairs and the other about the birth-
582 death pairs after the cancellation. We re-state and prove them separately in this appendix.

583 \triangleright **Claim A.1 (Impact on Shallow Pairs).** Let $\phi: X \rightarrow \mathbb{R}$ be a filter on a Lefschetz complex,
584 (s, t) a shallow pair of ϕ , and $\phi': X' \rightarrow \mathbb{R}$ the filter on the quotient after canceling (s, t) .
585 Then $\text{SH}(\phi) \subseteq \text{SH}(\phi') \cup \{(s, t)\}$.

586 **Proof.** Let $(u, v) \in \text{SH}(\phi)$ different from (s, t) . By Lemma 5.4 and the injectivity of the
587 map $\text{birth}: X_\times \rightarrow X_*$, u is different from s and t and so is v . Thus, u is still a facet of v in
588 X' . It is possible that v inherits additional facets from t , but only if s is a facet of v , in
589 which case all these new facets precede s and therefore also u in the linear ordering induced
590 by ϕ . Symmetrically, u may inherit additional cells it is a facet of from s , but only if u is
591 a facet of t , in which case all these new cells succeed t and therefore also v in the linear
592 ordering induced by ϕ . Restricting this ordering to X' , we get the linear ordering induced by
593 ϕ' , which implies that u is still the last facet of v , and v is still the first cell u is a facet of.
594 Equivalently, $(u, v) \in \text{SH}(\phi')$. ◀

595 \triangleright **Claim A.2 (Impact on Birth-death Pairs).** Let $\phi: X \rightarrow \mathbb{R}$ be a filter on a Lefschetz complex,
596 (s, t) a shallow pair of ϕ , and $\phi': X' \rightarrow \mathbb{R}$ the filter on the quotient after canceling (s, t) .
597 Then $\text{BD}(\phi) = \text{BD}(\phi') \cup \{(s, t)\}$.

598 **Proof.** We write y_1, y_2, \dots, y_n for the linear ordering of the cells in X induced by ϕ . Let
599 $k < \ell$ be the indices such that $s = y_k$ and $t = y_\ell$, and recall that removing s and t from this
600 list gives the linear ordering of the cells in $X' = X \setminus \{s, t\}$ induced by ϕ' . Let $c \in C(X)$ and
601 recall that $\pi(c) \in C(X')$ is obtained by dropping t and substituting the cells in $\partial t \setminus \{s\}$ for s ;
602 see the definition of this homomorphism in (8). The transition from c to $\pi(c)$ is implicit in
603 the operation that turns the Lefschetz complex of X into the quotient, which is the Lefschetz
604 complex of X' ; see Definition 3.1. Since (s, t) is a shallow pair, s is the last facet of t in the
605 linear ordering or, equivalently, the other facets of t precede s . It follows that the last cell of
606 c is also the last cell of $\pi(c) \cup \{s, t\}$ in the ordering. Hence,

$$607 \quad c \in C(X_b) \implies \pi(c) \in C(X'_b) \tag{23}$$

608 for any value $b \in \mathbb{R}$, in which we recall that $X_b \subseteq X$ and $X'_b \subseteq X'$ consist of all cells with
609 value at most b . Following the original algorithm for computing persistent homology in [8],
610 the remainder of this proof is inductive and establishes four hypotheses simultaneously. To
611 formulate them, let $b = \phi(y_j)$ and write d_j, d'_j for the canonical cycles associated with y_j ,
612 and $Y_j \subseteq X_* \cap X_b, Y'_j \subseteq X'_* \cap X'_b$ for the subsets of birth-giving cells such that the homology
613 classes of their canonical cycles form bases of $H(X_b)$ and $H(X'_b)$, respectively; see Lemma 5.2.
614 The hypotheses are

- 615 A. $d'_j = \pi(d_j)$ whenever $y_j \in X'_*$;
- 616 B. $y_j \in X_* \implies y_j \in X'_*$ and $y_j \in X_\times \implies y_j \in X'_\times$ whenever $y_j \in X'$;
- 617 C. $Y'_j = Y_j$ whenever $j < k$ or $\ell < j$, and $Y'_j = Y_j \setminus \{s\}$ whenever $k < j < \ell$;
- 618 D. $\text{birth}'(y_j) = \text{birth}(y_j)$ whenever $y_j \in X'_\times$.

619 By construction, $d_j \setminus \{y_j\} \subseteq X_\times$, which implies $s \notin d_j$ unless $s = y_j$. Hence, Hypothesis A is
620 equivalent to $d'_j = d_j \setminus \{t\}$. Hypothesis B is equivalent to $X'_* = X_* \cap X'$ and $X'_\times = X_\times \cap X'$.
621 By Lemma 5.4, we have $s \in Y_j$ iff $k \leq i < \ell$, which implies that Hypothesis C is equivalent

XX:18 The Depth Poset of a Filtered Lefschetz Complex

622 to $Y'_j = Y_j \setminus \{s\}$ for all indices $j \neq k, \ell$. Hypothesis D readily implies that with the exception
 623 of (s, t) , the birth-death pairs are the same before and after the cancellation.

624 The four hypotheses are void and thus trivially true for $j = 0$, which serves as the
 625 induction basis. For the inductive step, consider a cell y_j , with value $b = \phi(y_j)$, and assume
 626 the four hypotheses are true for all indices $i \leq j - 1$. Let $a = \phi(y_{j-1})$, and assume $y_j \neq s, t$,
 627 else there is nothing to prove.

628 Consider first the case $y_j \in X_*$, so the canonical cycle associated to y_j , denoted $d_j \subseteq$
 629 $(X_x \cap X_a) \cup \{y_j\}$, is well defined. Let $d'_j = \pi(d_j)$. By construction, y_j belongs to d'_j , and by
 630 (23) and the inductive assumption, all other cells in d'_j belong to $X'_x \cap X'_a$. Indeed, $s \notin d_j$
 631 because it gives birth, so $d'_j \subseteq d_j$ because the cancellation of (s, t) does not add any new cells
 632 to the cycle. Hence, d'_j is the canonical cycle of y_j in X' , which establishes Hypothesis A.
 633 But this also shows $y_j \in X_* \Rightarrow y_j \in X'_*$, which is Hypothesis B for birth-giving cells. In
 634 addition, it shows $Y_j = Y_{j-1} \cup \{y_j\}$ and $Y'_j = Y'_{j-1} \cup \{y_j\}$, which together with the inductive
 635 assumption implies Hypothesis C for birth-giving cells.

636 Consider second the case $y_j \in X_x$. Then ∂y_j is a non-trivial cycle in X_a . By (23),
 637 $\partial' y_j = \pi(\partial y_j)$ is a cycle in X'_a . To establish that it is non-trivial, assume there is a chain,
 638 $c_j \subseteq X'_a$ with $\partial' c_j = \partial' y_j$. We distinguish a few cases and conclude that ∂y_j is trivial in
 639 each, which is a contradiction to the assumption and thus implies that $\partial' y_j$ is non-trivial.

- 640 1. $j < k$. Then $\partial' y_j = \partial y_j$ and $\partial' c_j = \partial c_j$. Hence, ∂y_j is trivial, contradiction.
- 641 2. $k < j < \ell$. Then there cannot be any cycle homologous to ∂y_j that contains s . Indeed, if
 642 there is such a cycle, then s is a facet of a cell in X_b , which contradicts that t is the first
 643 cell that has s as a facet in the linear ordering. Hence, $\partial' y_j = \partial y_j$, and since $c_j \subseteq X_a$, we
 644 also have $\partial c_j = \partial y_j$, so again ∂y_j is trivial, contradiction.
- 645 3. $\ell < j$. If $s, t \notin \partial y_j$, then we use the same argument as above to show that ∂y_j is trivial,
 646 contradiction. This leaves two subcases.
 - 647 3.1 $s \in \partial y_j$. Then $\partial' y_j = \partial y_j + \partial t$. Hence, $\partial y_j = \partial' c_j + \partial t = \partial(c_j \cup \{t\})$, which implies
 648 that ∂y_j is trivial, contradiction.
 - 649 3.2 $t \in \partial y_j$. Then $\partial' y_j = \partial y_j \setminus \{t\}$, and similarly, $\partial' c_j = \partial c_j \setminus \{t\}$. But then $\partial y_j = \partial c_j$, so
 650 ∂y_j is trivial, contradiction.

651 In all cases, we get $y_j \in X'_x$, which establishes Hypothesis B for death-giving cells. Next, we
 652 show that y_j is paired with the same birth-giving cell before and after the cancellation; that
 653 is: $\text{birth}'(y_j) = \text{birth}(y_j)$. This will establish Hypothesis D, which then together with the
 654 inductive assumption establishes Hypothesis C. Let $A \subseteq Y_j$ be the birth-giving cells such
 655 that $d = \sum_{x \in A} d_x$ satisfies $[d]_a = [\partial y_j]_a$, and similarly let $A' \subseteq Y'_j$ be the birth-giving cells
 656 such that $d' = \sum_{x \in A'} d'_x$ satisfies $[d']_a = [\partial' y_j]_a$. In Case 1, we have $Y'_j = Y_j$ and $d'_i = d_i$
 657 for every $y_i \in Y'_j$, so $A' = A$. In Case 2, we have $Y'_j = Y_j \setminus \{s\}$, but as argued there, $s \notin A$,
 658 which again implies $A' = A$. In Case 3, we have $Y'_j = Y_j$ and $d'_i = d_i \setminus \{A\}$ for every $y_i \in Y'_j$,
 659 so we get $A' = A$ in all subcases. We thus get $\text{birth}'(y_j) = \text{birth}(y_j)$ in all three cases.

660 This completes the inductive argument, which implies $\text{BD}(\phi') = \text{BD}(\phi) \setminus \{(s, t)\}$, as
 661 required to establish the claimed relation. \blacktriangleleft

662

B Notation

663

 $(X, \dim, \kappa), (Y, \dim|_Y, \kappa|_{Y \times Y})$ Lefschetz complex, subcomplex

664

 $\partial: C(X) \rightarrow C(X)$ boundary homomorphism

665

 $C(X), Z(X), B(X), H(X)$ chains, cycles, boundaries, homologies

666

 c, d chains

667

668

 (X', \dim', κ') quotient Lefschetz complex

669

 $s, t; u, v$ reducible pairs

670

 w, x, y, z cells

671

 V, G_V combinatorial gradient, associated digraph

672

 x_0, x_1, \dots, x_n path, connection, cycle

673

 $f: X \rightarrow \mathbb{R}$ Lyapunov function

674

675

 $\phi: X \rightarrow \mathbb{R}$ filter

676

 X_*, X_\times birth-giving, death-giving cells

677

 $X_b = \phi^{-1}(-\infty, b]$ sublevel set

678

 $Y'_b \subseteq Y_b$ birth-giving cells alive at b

679

 $b = \phi(y); d_y; [d_y]_b$ value; canonical cycle; homology class

680

681

birth: $X_\times \rightarrow X_*$ birth function

682

 $\text{BD}(\phi) = \{(\text{birth}(y), y)\}$ birth-death pairs

683

 $\text{SH}(\phi)$ shallow pairs

684

 $\text{Depth}(\phi)$ depth poset

685

 $S, S_{(u,v)} \subseteq \text{SH}(\phi_{j-1})$ subsets of shallow pairs

686

 $\phi_j: X_j \rightarrow \mathbb{R}, \phi_j^S: X_j^S \rightarrow \mathbb{R}$ j -th derived filter, after canceling pairs

■ Table 1: Notation used in the paper.

687 **C** Results and Definitions

- 688 ■ Section 1: Introduction.
- 689 ■ Section 2: Lefschetz Complexes.
 - 690 ■ Definition 2.1 (Lefschetz Complex).
- 691 ■ Section 3: Cancelling Reducible Pairs.
 - 692 ■ Definition 3.1 (Cancellation).
 - 693 ■ Lemma 3.2 (Quotient).
 - 694 ■ Lemma 3.3 (Chain Homotopy).
- 695 ■ Section 4: Vectors of Combinatorial Gradients.
 - 696 ■ Lemma 4.1 (Lyapunov Function).
 - 697 ■ Lemma 4.2 (Independence and Acyclicity).
 - 698 ■ Theorem 4.3 (Morse Complex).
- 699 ■ Section 5: Shallow Pairs in Persistent Homology.
 - 700 ■ Lemma 5.1 (Canonical Cycle Basis).
 - 701 ■ Lemma 5.2 (Canonical Homology Basis).
 - 702 ■ Definition 5.3 (Shallow Pairs).
 - 703 ■ Lemma 5.4 (Shallow Pairs are Special).
 - 704 ■ Theorem 5.5 (Canceling a Shallow Pair).
- 705 ■ Section 6: The Depth Poset.
 - 706 ■ Lemma 6.1 (Shallow Pairs as Vectors).
 - 707 ■ Lemma 6.2 (Turning Shallow).
 - 708 ■ Definition 6.3 (Depth Poset).
 - 709 ■ Theorem 6.4 (Cancelable Linear Orders).
 - 710 ■ Theorem 6.5 (Ordering by Persistence).
- 711 ■ Section 7: Discussion.
- 712 ■ Appendix A: Proof of Theorem 5.5.
 - 713 ■ Claim A.1 (Impact on Shallow Pairs).
 - 714 ■ Claim A.2 (Impact on Birth-death Pairs).

715 **D To Think and to Do**

- 716 ■ Section 1: Introduction.
- 717 ■ Section 2: Lefschetz Complexes.
- 718 ■ Section 3: Cancelling Reducible Pairs.
- 719 ■ Section 4: Vectors of Combinatorial Gradients.
- 720 ■ Section 5: Shallow Pairs in Persistent Homology.
- 721 ■ Section 6: The Depth Poset.
- 722 ■ Section 7: Discussion.
- 723 ■ Appendix A: Proof of Theorem 5.5.